

Distributed Admission Control for Voice over IP

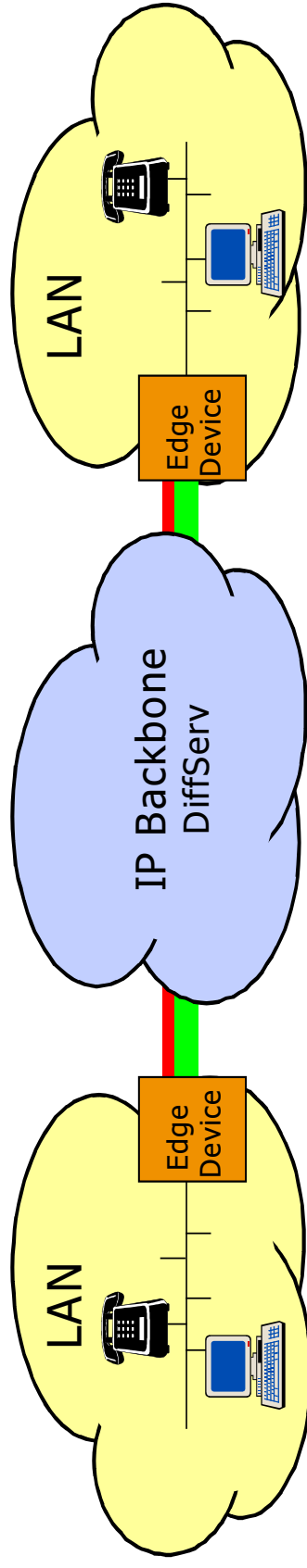
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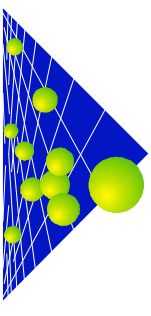
ITG Workshop

IP Netzplanung, IP Netzmanagement und Optimierung

Universität Würzburg, 20.-21. Juli 2000

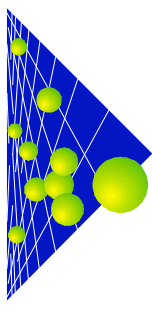


Voice over IP

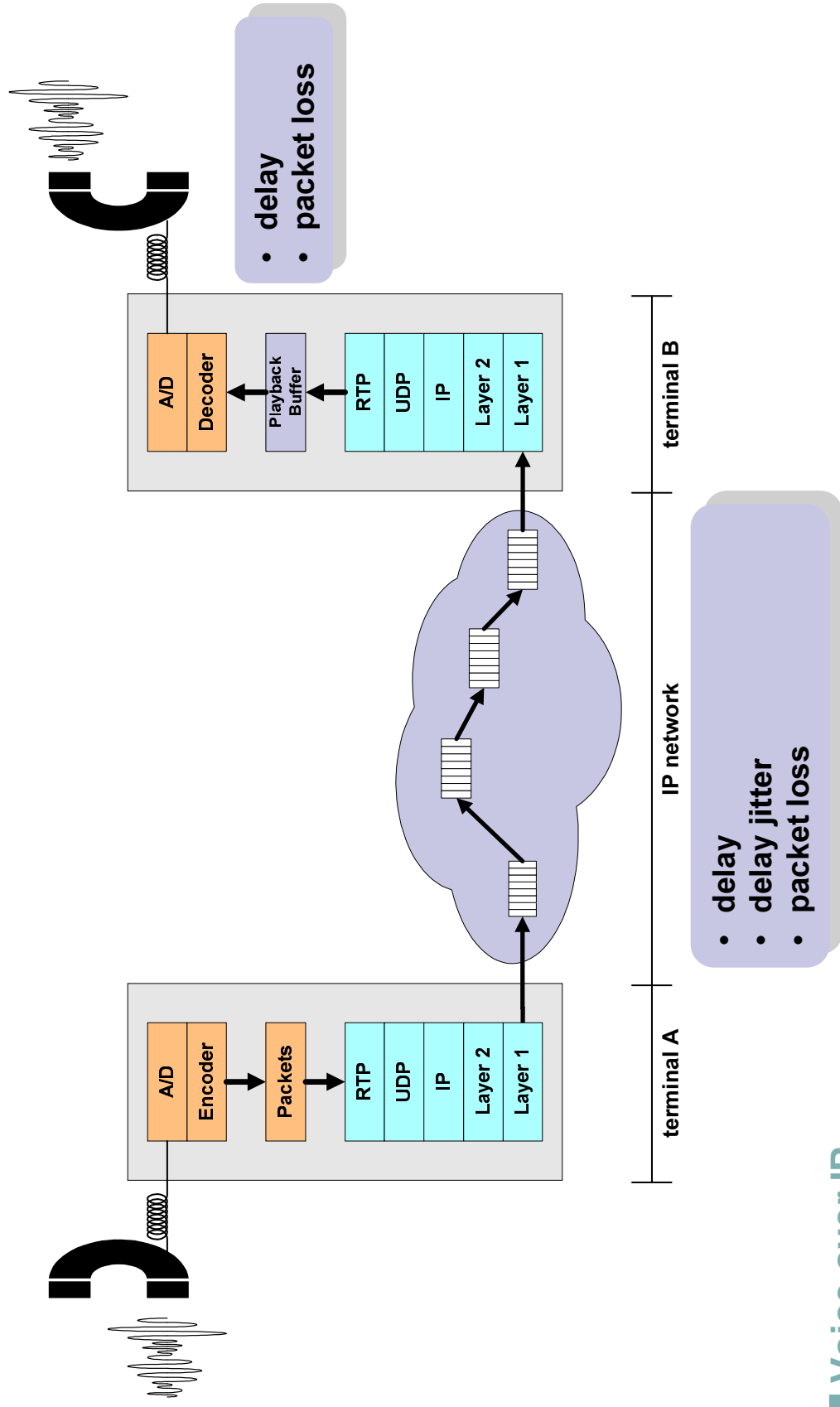


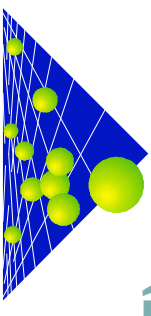
Outline

- ◆ Voice over IP
- ◆ Admission Control
- ◆ Dimensioning
- ◆ Simulation Model
- ◆ Simulation Results
- ◆ Conclusions

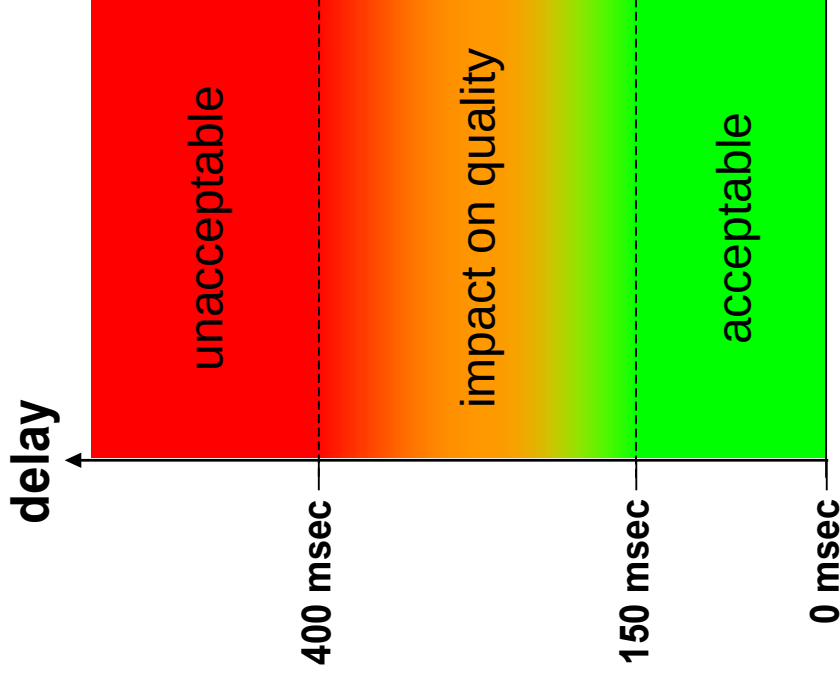


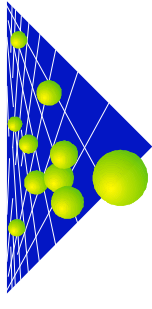
Voice over IP and QoS Problems





G.114: One-Way Transmission Time (ITU-T)

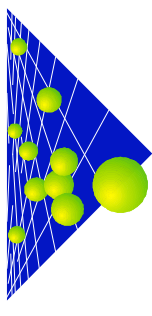




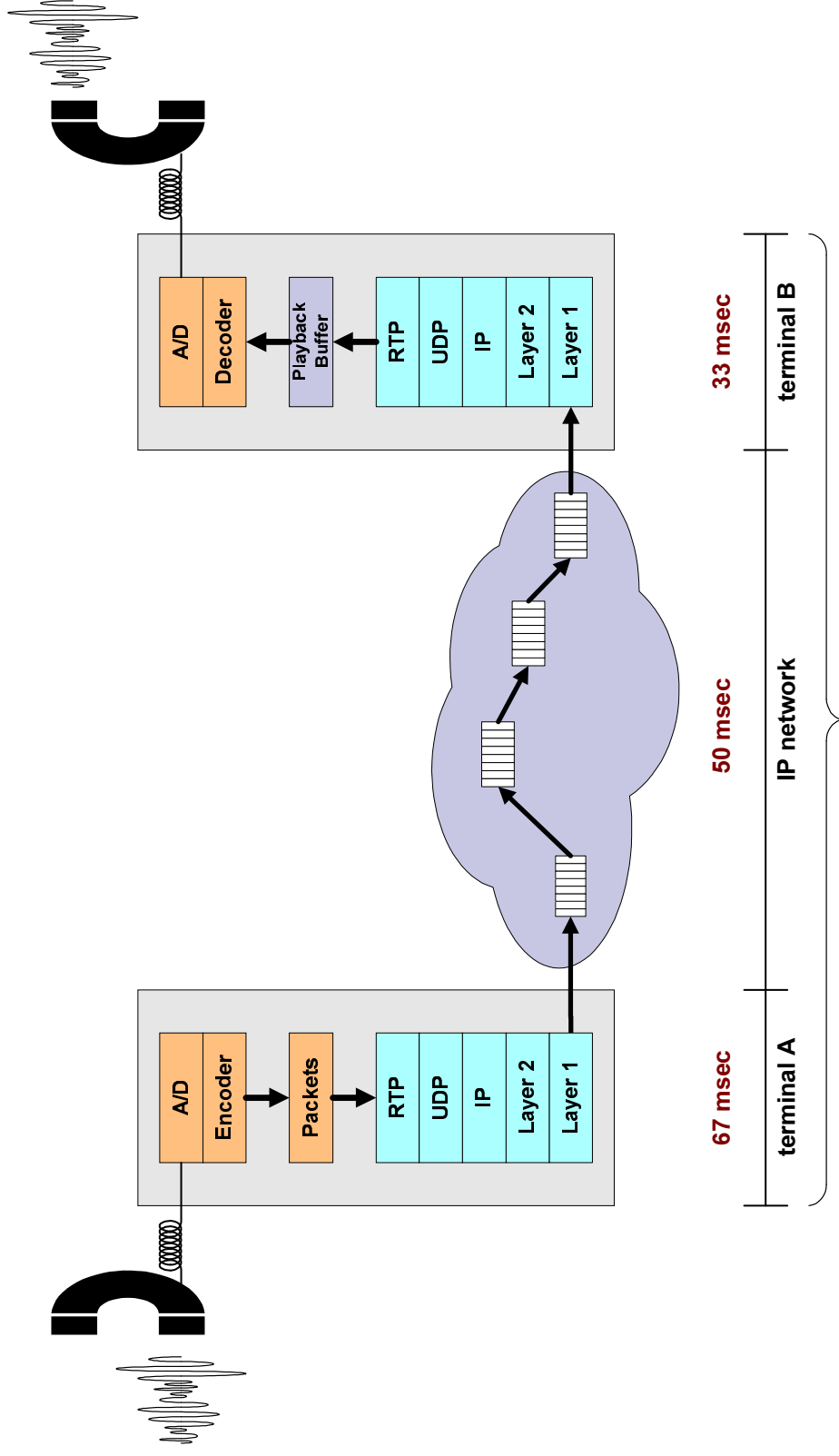
Voice Codecs

	G.723.1	G.729
Delay Encoder / msec		
• frame size	30,0	10
• lookahead	7,5	5
• processing	30,0	10
Delay Decoder / msec	32,5	12
Σ delay / msec	100,0	37

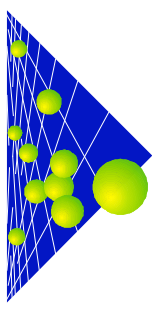
Data Rate Code / Kbps	5,3 / 6,4	8
Frame Size / Byte	20 / 24	10
Overhead (IP, UDP, RTP) / Byte	40	40
IP data rate / Kbps	15,9 / 17	40



VoIP

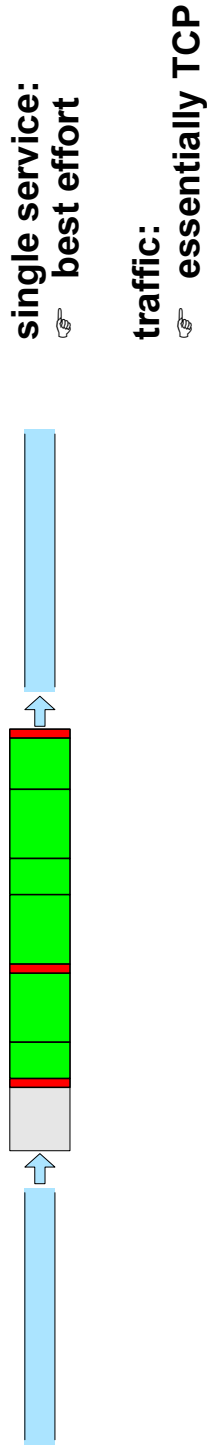


ITU, G.114: end-to-end delay < 150 msec

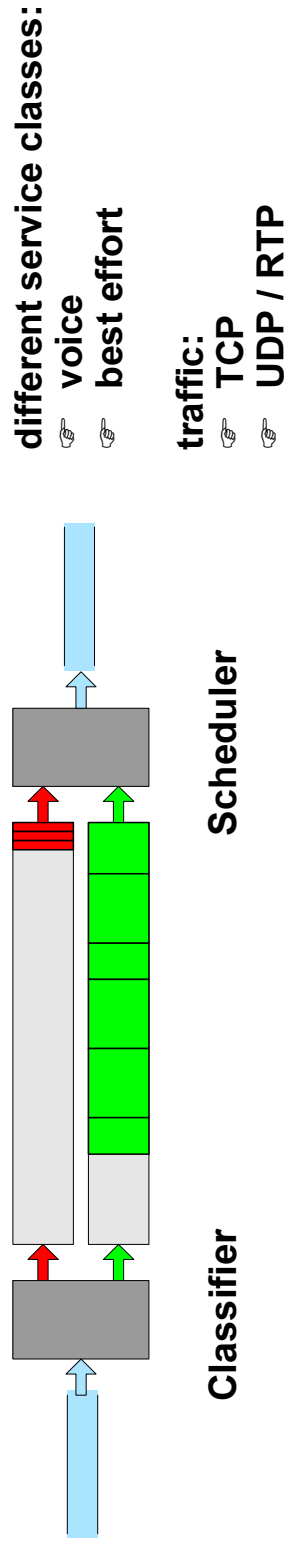


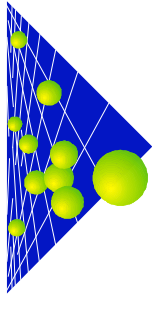
Differentiated Services (DiffServ)

IP today:

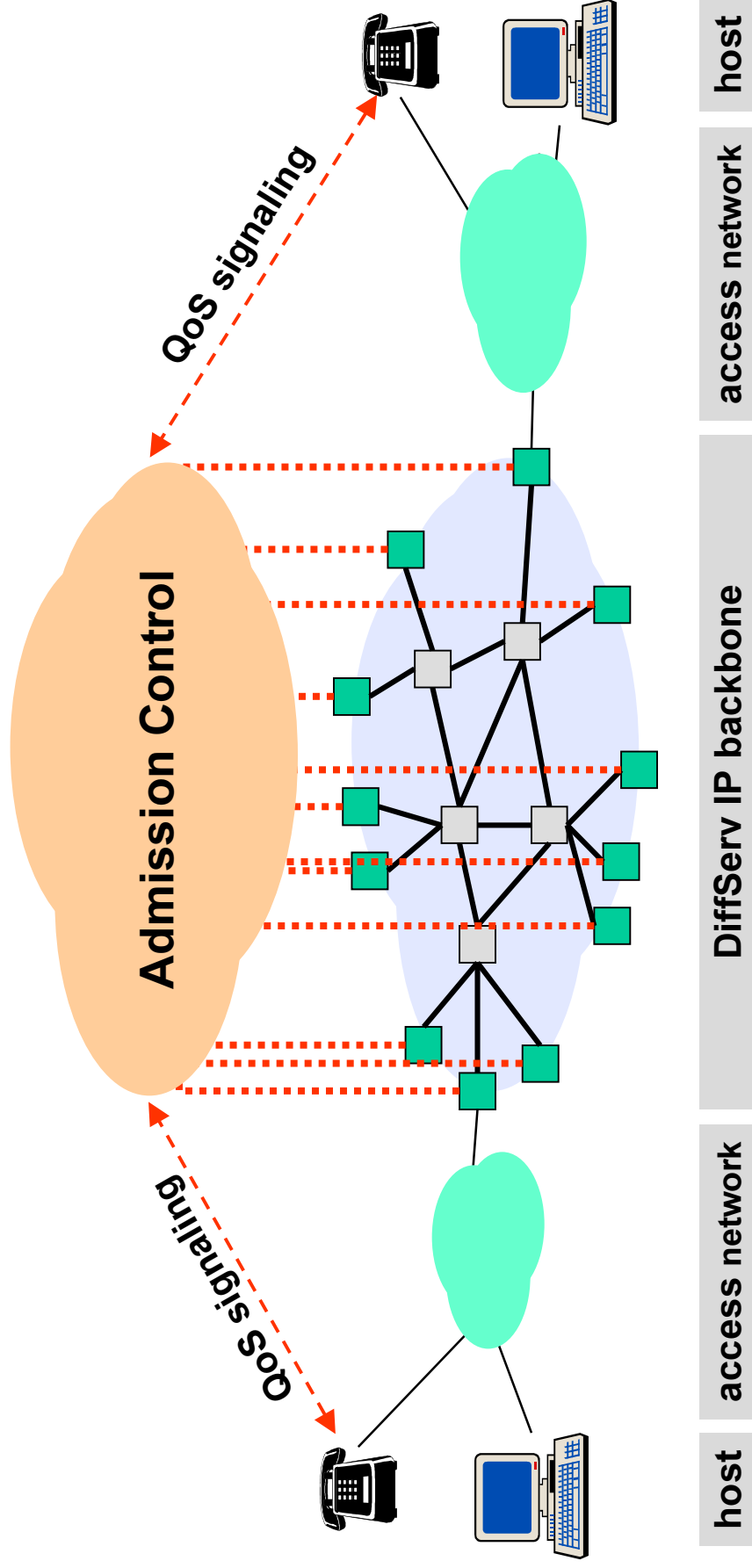


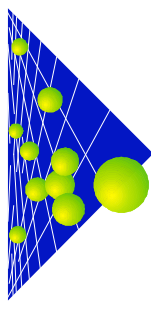
DiffServ:



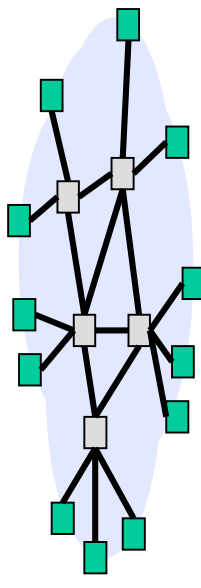
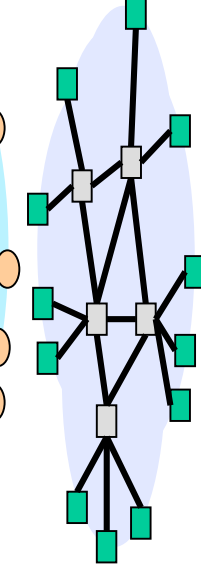
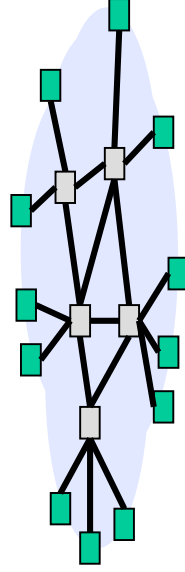
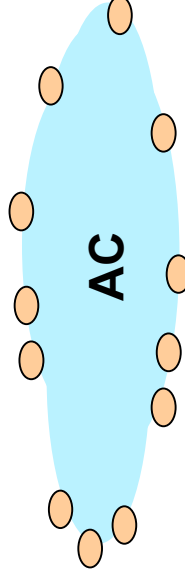
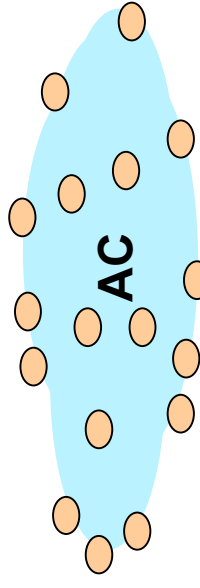
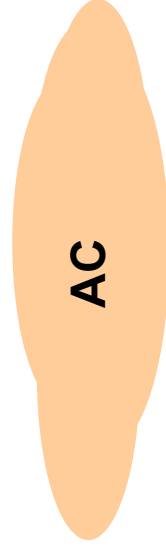


Network Scenario





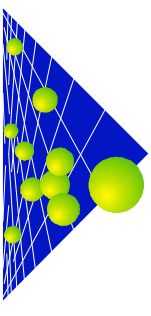
Admission Control



central

at network border only

hop-by-hop



Admission Control

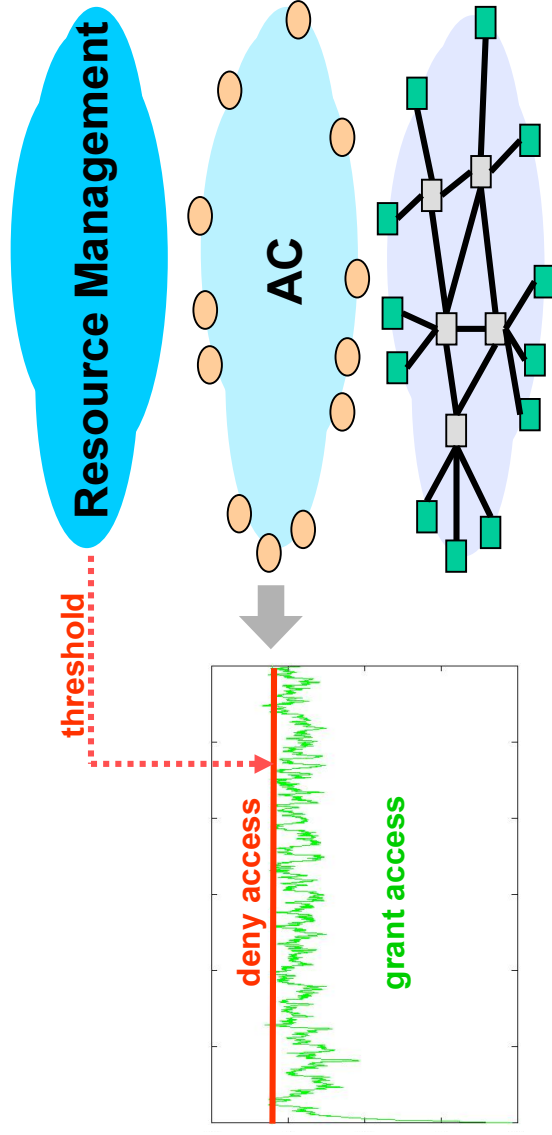
◆ local bandwidth table:

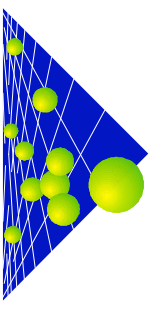
- ☞ available bandwidth for each egress edge device: C_E
 - provided by resource management
 - will be adapted to changing conditions by resource management
- Erlang formula
 - traffic matrix
 - $B = 0.01$

☞ bandwidth in use: A_E

◆ local decision:

☞ if $(A_E + a_E \leq C_E)$
 then grant access
 else deny access





Dimensioning

Method A

- ◆ calculate bandwidth demand for each end-to-end traffic stream
 - ☞ end-to-end traffic matrix
 - ☞ Erlang formula with $B=0.01$
- ◆ accumulate bandwidth demand for each link

conservative

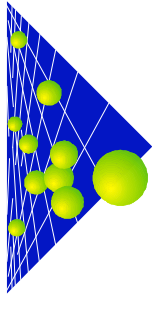
overprovisioning

Method B

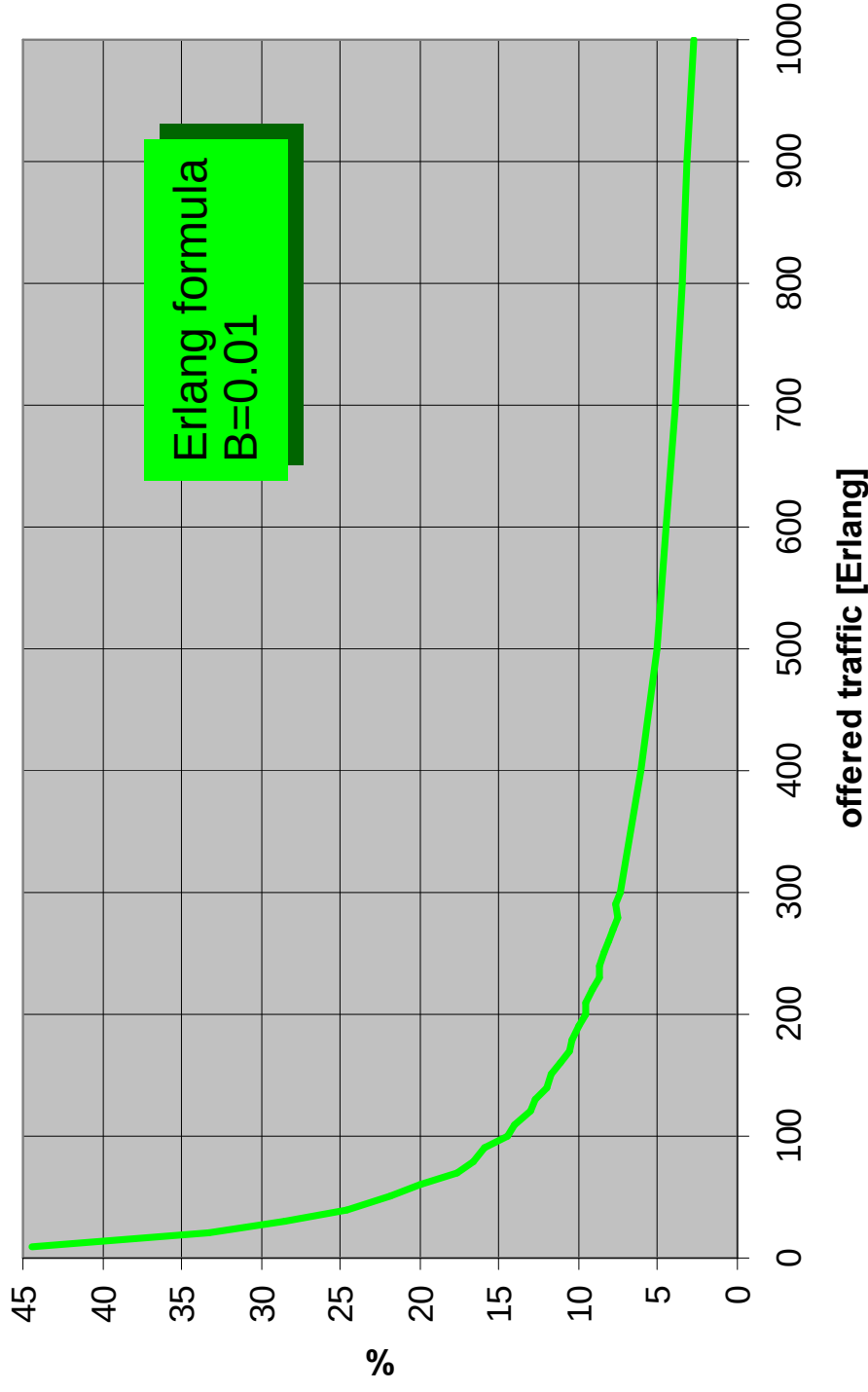
- ◆ accumulate offered traffic for each link
 - ☞ end-to-end traffic matrix
- ◆ calculate bandwidth demand of each single link
 - ☞ Erlang formula with $B=0.01$

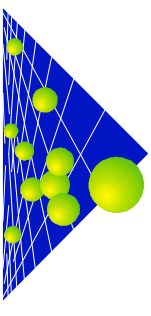
multiplex gains

probes level
of resource
optimization



An Upper Limit for Multiplex Gain of Dimensioning Method B





Voice Traffic Model

☞ Voice Codec

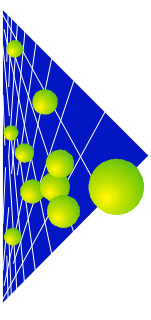
- ☞ G.723.1

☞ IP Packets

- ☞ 3 voice frames per packet
- ☞ 112 Byte
- ☞ 8 kbps CBR

☞ Voice Activity

- ☞ Poisson Model
- ☞ exponential holding times, mean 180 sec
- ☞ exponential inter-arrival times



Data Traffic Model

☞ self-similar traffic

☞ TCP

☞ Pareto distributed file sizes

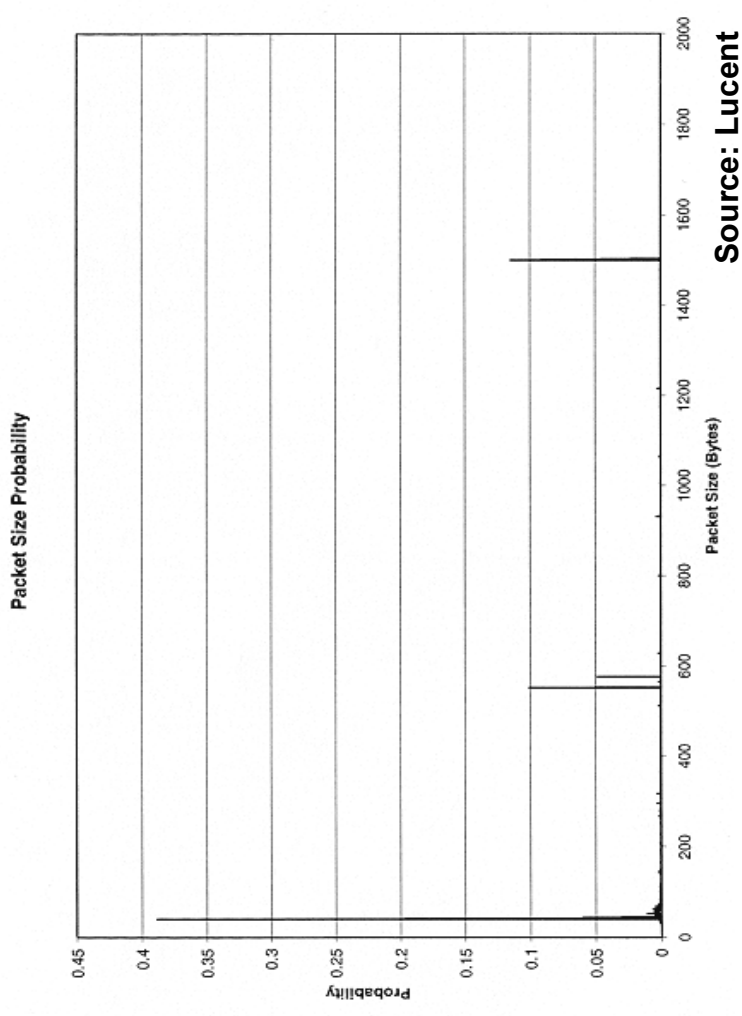
☞ exponential distributed inter arrival times

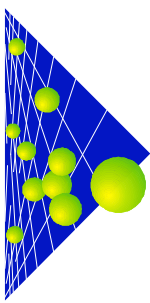
☞ 3 packet sizes [Byte]:

☞ 40 (acknowledgements)

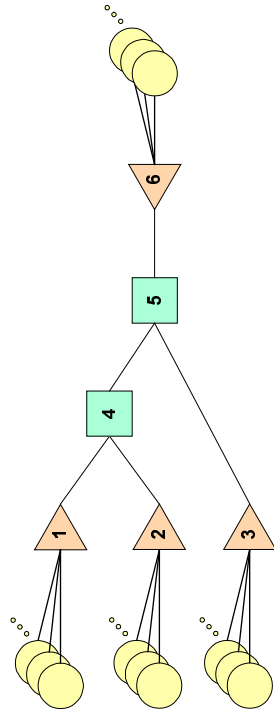
☞ 576 (default MTU size)

☞ 1500 (Ethernet)





Scenario A: Voice over IP

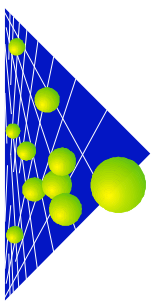


Traffic Load		
link	voice users	mean data rate [kbps]
1 → 6	240	240
2 → 6	480	480
3 → 6	400	400

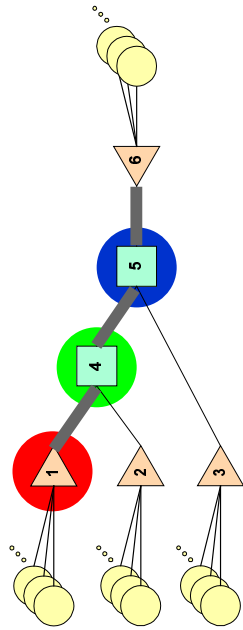
Queue Size
50 Packets

link	Network Dimensioning				
	A		B		
	rate [kbps]	utilisation [%]	rate [kbps]	utilisation [%]	saving [%]
1 → 4	336	71	336	76	0
2 → 4	600	80	600	80	0
3 → 5	512	78	512	79	0
4 → 5	936	77	856	84	8,5
5 → 6	1448	77	1272	88	12,2

Gatekeeper Thresholds		
GK	A	B
1	42	42
2	75	75
3	64	64
6	181	159

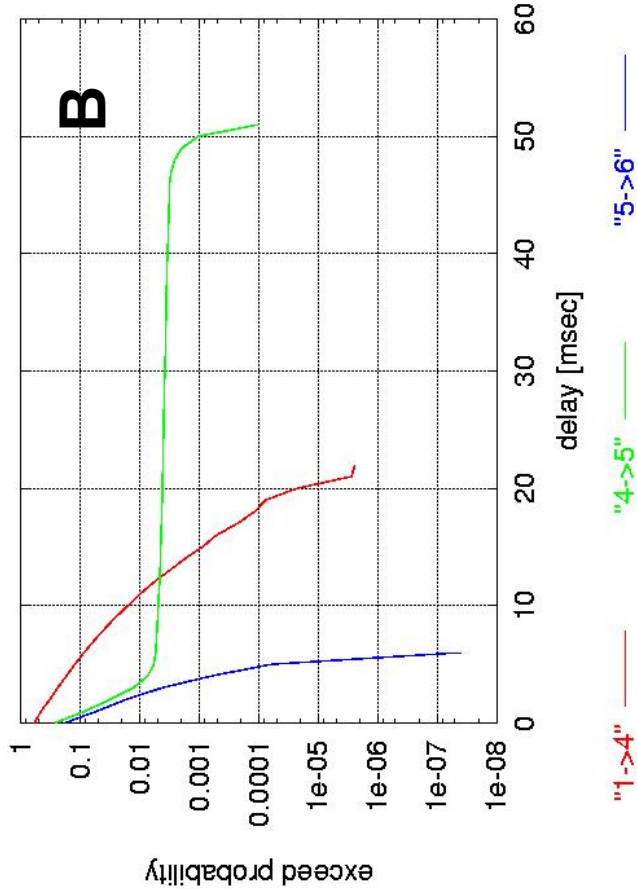


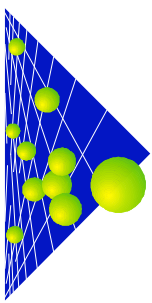
Scenario A: Delay, Packet Loss Rate



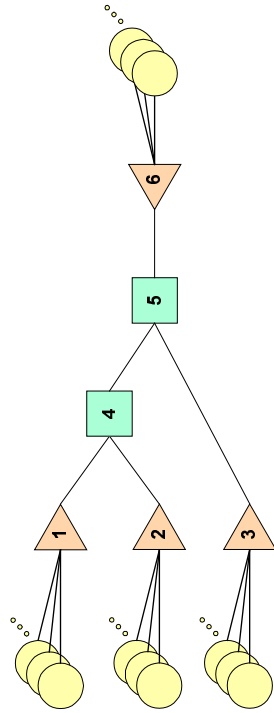
Mean Delay [msec]		
link	A	B
1→4	3,4	2,4
4→5	0,3	0,9
5→6	0,2	0,6

Packet Loss Frequency		
link	A	B
1→4	0	0
4→5	0	$4,4 \cdot 10^{-5}$
5→6	0	0





Scenario B: Voice + Data Traffic

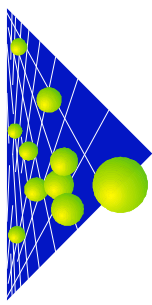


Traffic Load			
link	voice users	mean data rate	
		voice [kbps]	data [kbps]
1 → 6	240	240	270
2 → 6	480	480	480
3 → 6	400	400	410

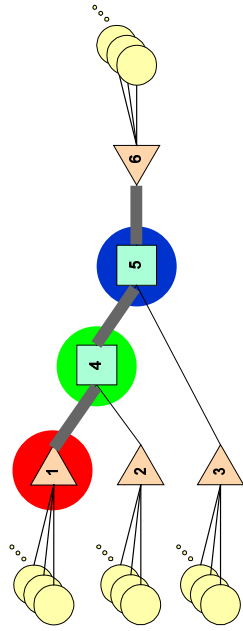
Queue Size
50 Packets

link	Network Dimensioning				
	A		B		
	rate [kbps]	utilisation [%]	rate [kbps]	utilisation [%]	saving [%]
1 → 4	672	76	672	76	0
2 → 4	1200	80	1200	80	0
3 → 5	1024	79	1024	79	0
4 → 5	1872	79	1792	82	4,3
5 → 6	2896	79	2720	84	6,1

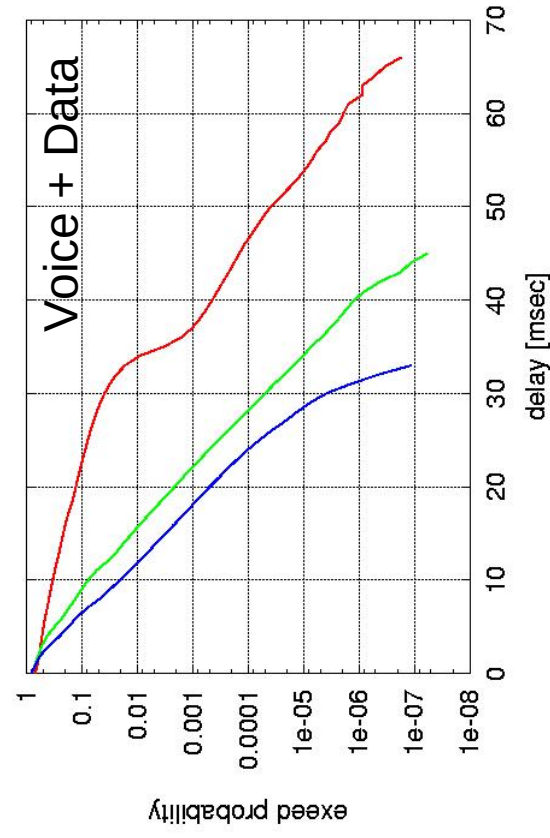
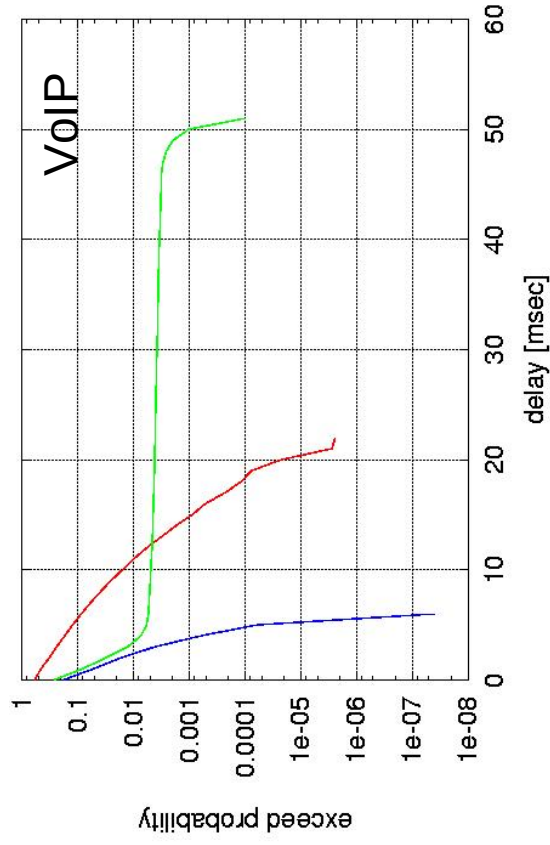
Gatekeeper Thresholds		
GK	A	B
1	42	42
2	75	75
3	64	64
6	181	159

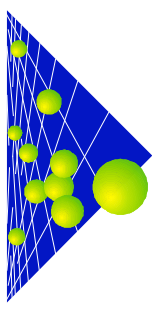


Influence of Data Traffic

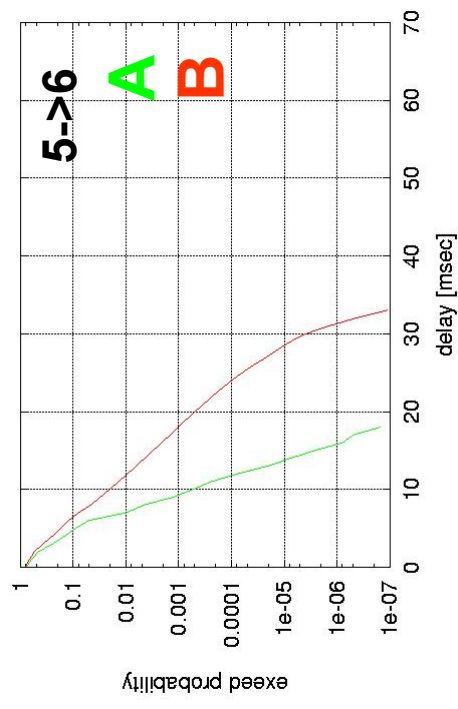
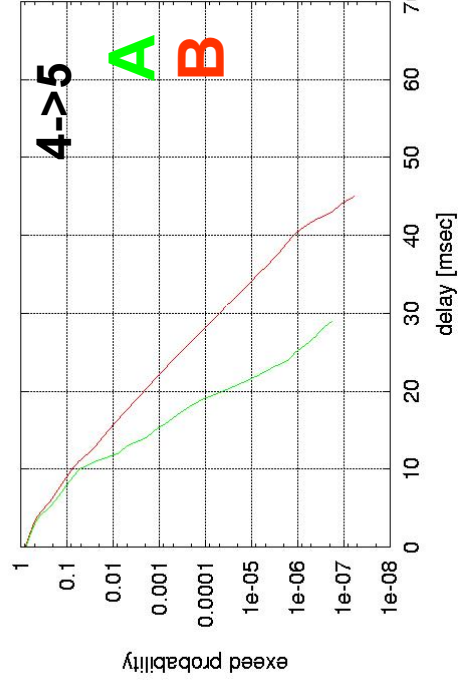
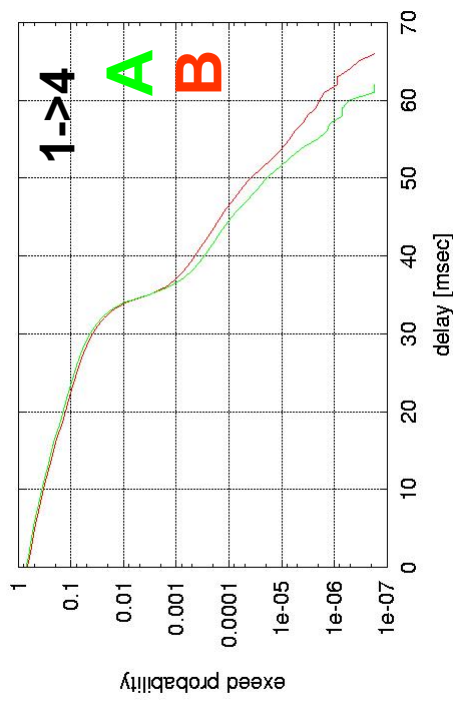
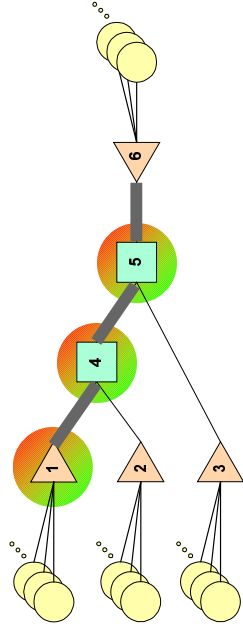


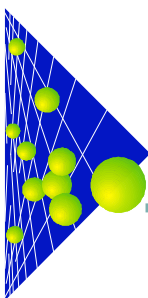
link	Mean Delay [msec]		Packet Loss Rate	
	VoIP	IVDoIP	VoIP	IVDoIP
1→4	2,4	8,7	0	0
4→5	0,9	4,6	$4,4 \cdot 10^{-5}$	0
5→6	0,6	3,6	0	$1,5 \cdot 10^{-7}$



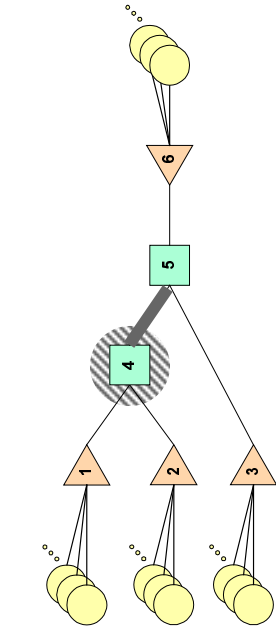


Comparison of Dimension Methods (1)

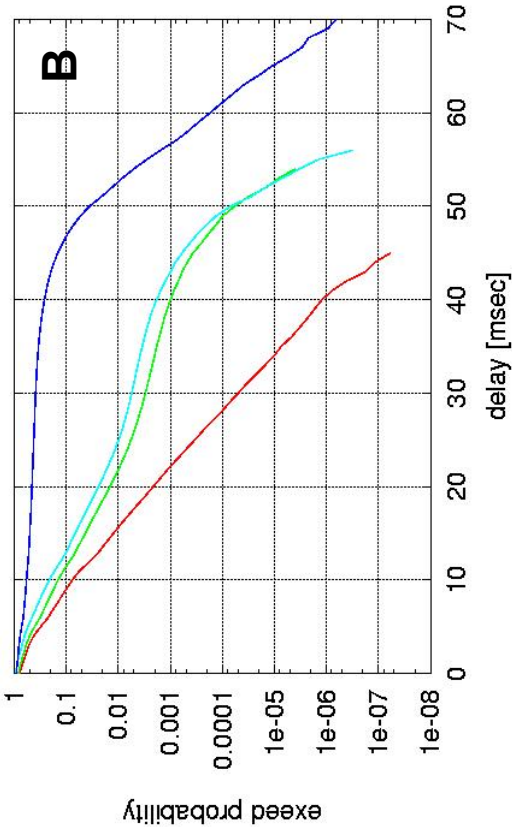
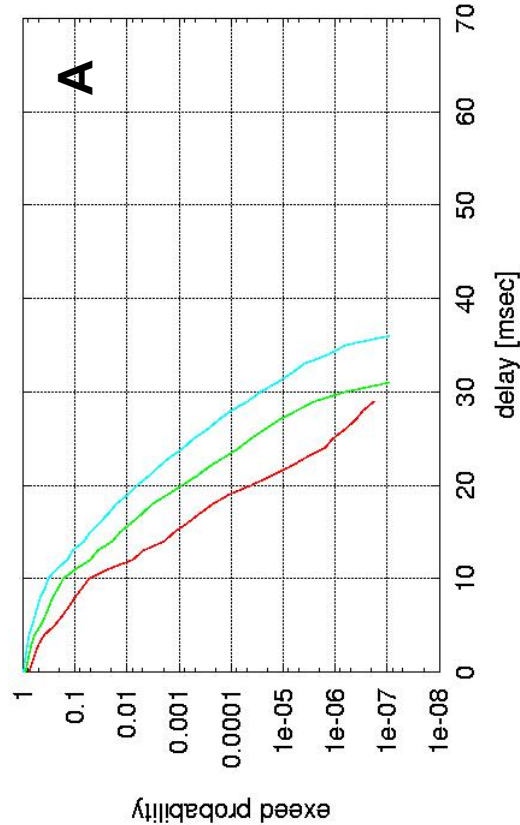




Comparison of Dimension Methods (2): Overload

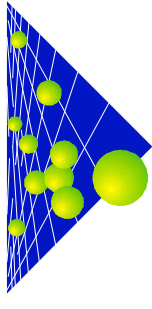


Scenario	Load			
	A		B	
	voice	data	voice	data
	79%	80%	82%	80%
	100%	80%	100%	80%
	100%	160%	100%	160%
			Schiefplast	80%

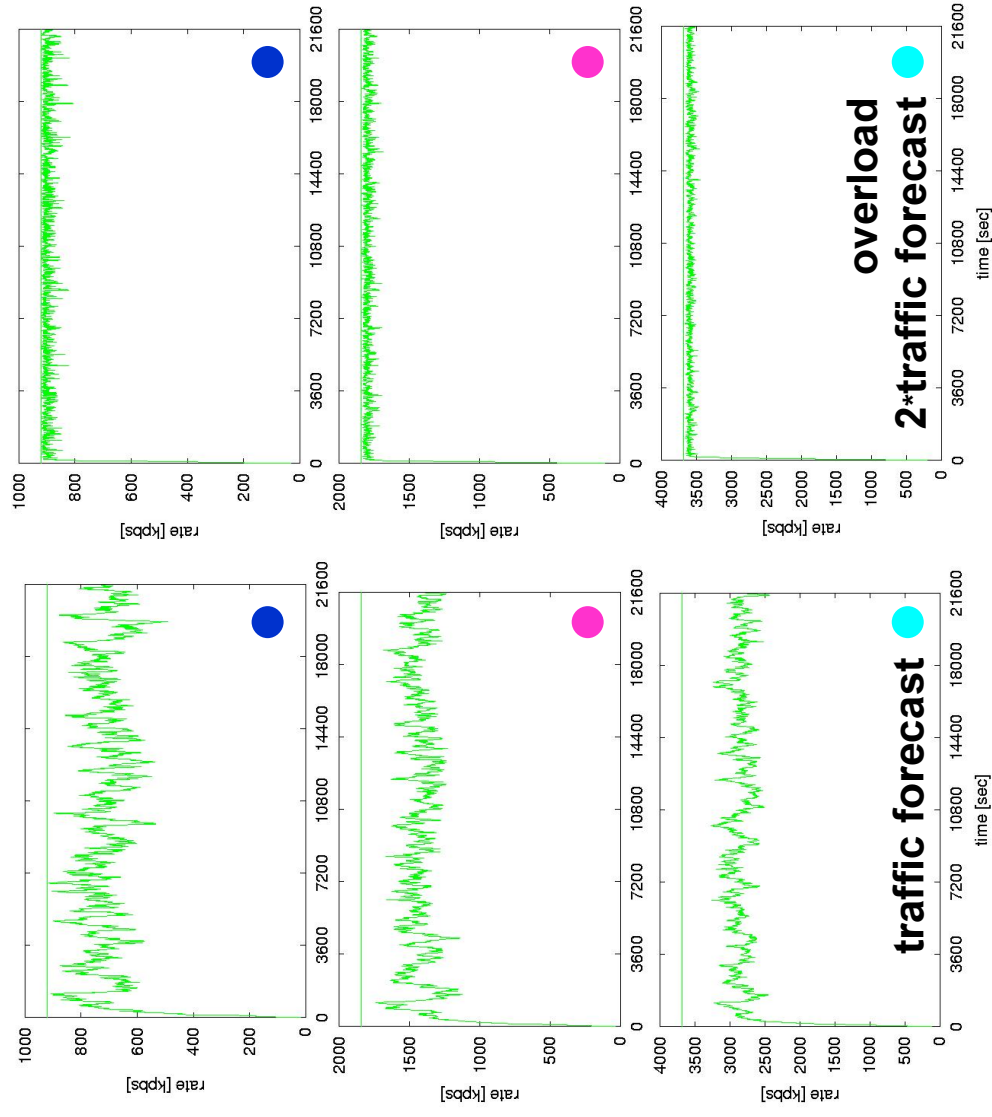
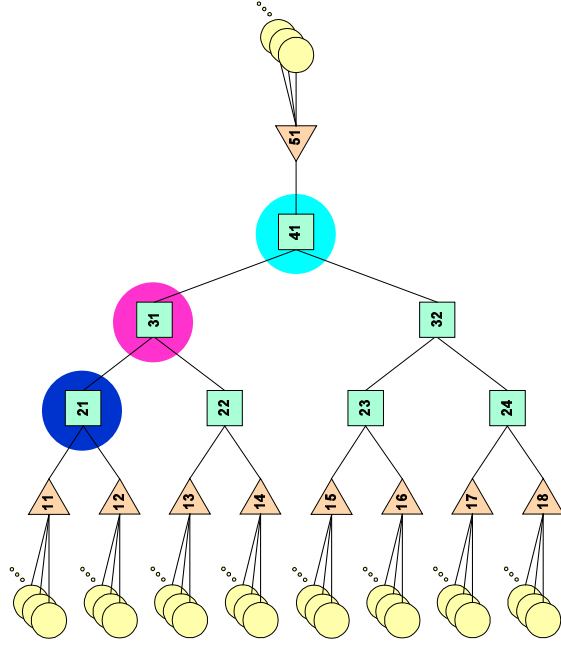


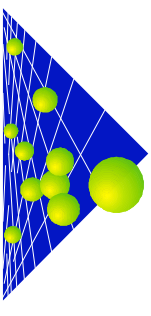


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Utilisation (2): Overload





Conclusions (1)

both traffic classes gain from integration

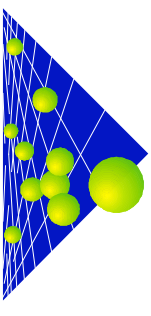
- ☞ they “borrow” unused resources

Dimensioning Method A

- ☞ good QoS for voice traffic
- ☞ QoS is kept in overload scenarios
- ☞ no congestion
- ☞ overprovisioning

Dimensioning Method B

- ☞ capacity saving
- ☞ congestion
- ☞ light QoS problems if traffic load meets traffic forecast (about 1% of packets are affected)
- ☞ heavy QoS problems in overload scenarios



Conclusions (2)

Independent Local Admission Control

- ☞ is feasible
- ☞ user perceived voice quality is very sensitive to resource management

Next Steps

- ☞ signalling for admission control
- ☞ tracking used bandwidth
- ☞ resource management
- ☞ advanced independent local decision schemes