

The efficiency of Traffic Engineering (TE) with Regard to Failure Resilience

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- Focus on capacity upgrading process for steady traffic increase
- Strategy for link upgrades with TE & failure resilience yielding maximum utilization
- Examples & Outlook

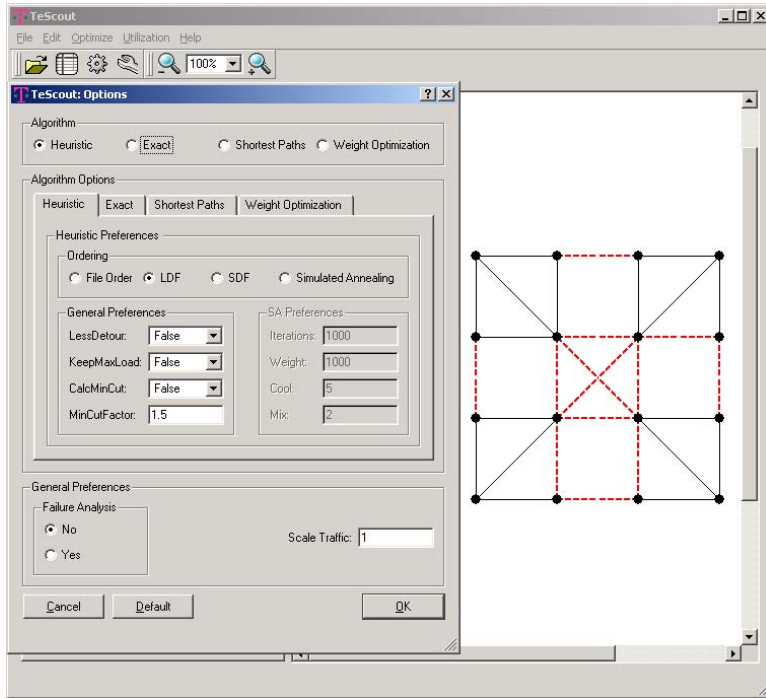
Traffic Engineering (TE) & Multiprotocol Label Switching

- Shortest path routing in pure IP networks without TE leads to unbalanced load, which cannot be efficiently controlled
- Multiprotocol label switching (MPLS) provides
 - Label Switched Paths (LSP) from source to destination
 - Measurement per LSP to obtain the traffic matrix
 - Full mesh of LSP between edge routers: difficult to operate
 - QoS support (for DiffServ & IntServ)
 - Virtual private networks (VPN)
- Integration of MPLS with the optical layer (GMPLS, ASON)

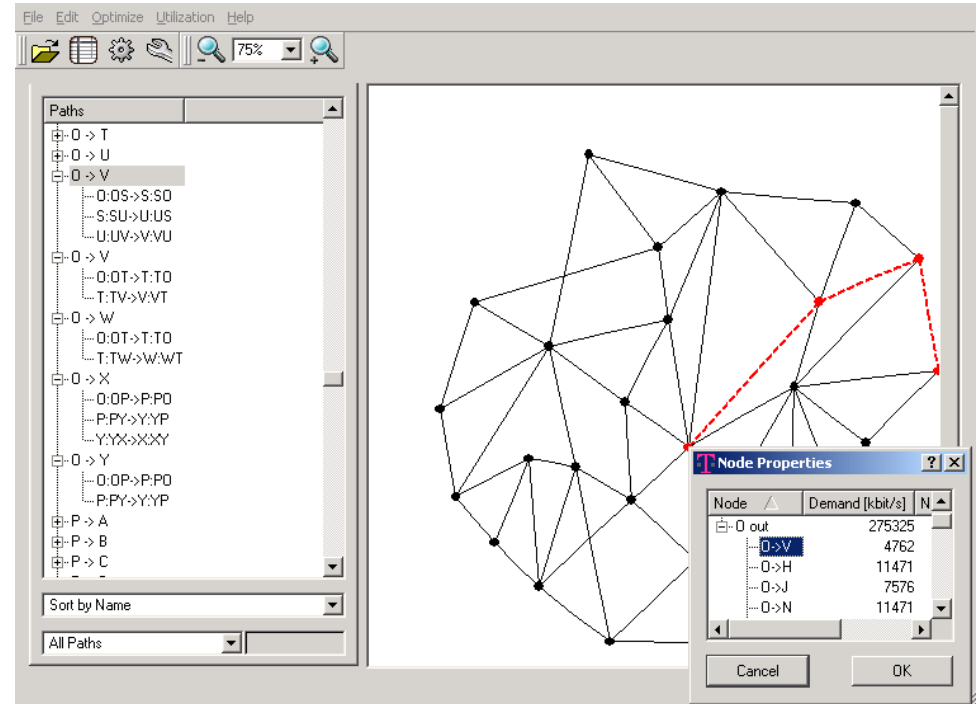
Traffic Engineering: Algorithms for Load-Balancing

- Main optimization goal
 - Minimize the maximum load on the network links ⇒ Allows for largest overall traffic increase until next upgrade when some load threshold is exceeded
- Optimization algorithms: *TE-Scout* tool developed at T-Systems
 - Simulated annealing (single path per demand)
 - Linear programming (multiple paths per demand possible)
 - Algorithm based on max-flow-min-cut principle
 - Link weight adaptationHeuristics for pure IP networks studied by [Fortz & Thoroup]

TE-Scout Tool: Optimization & representation of path designs

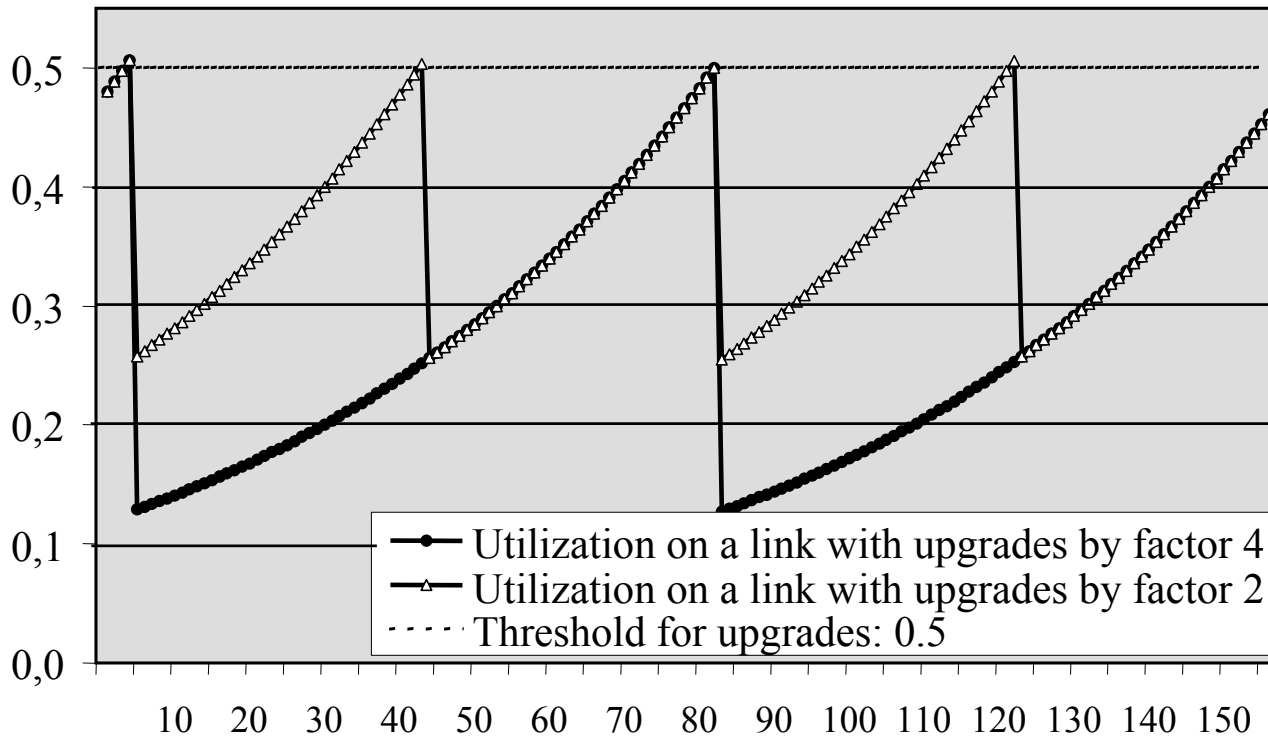


Options for applying optimization algorithms



Visualization of paths for traffic demands

Utilization gaps for link upgrades without TE



IP link upgrades to 2- (or 4-)fold capacity reduce the link load to $\frac{1}{2}$ ($\frac{1}{4}$)

Traffic engineering (TE) can use new capacity after a reoptimization with traffic shifts on link

TE yields more constant load close to threshold

Mean load u for threshold based upgrades without TE:

$$u = \frac{T_{\text{Upgrade}}}{F_{\text{Upgrade}} D_{\text{Upgrade}}} \int_0^{D_{\text{Upgrade}}} \exp(\omega x) dx = T_{\text{Upgrade}} \frac{1 - 1/F_{\text{Upgrade}}}{\ln(F_{\text{Upgrade}})} ; \quad \begin{aligned} F_{\text{Upgrade}} = 2 &\Rightarrow u = 0.72 T_{\text{Upgrade}} \\ F_{\text{Upgrade}} = 4 &\Rightarrow u = 0.54 T_{\text{Upgrade}} \end{aligned}$$

Capacity expansion with TE regarding failures

Assumption: link upgrades are done to twice the capacity

Expansion strategy to decide, which link upgrade to do next:

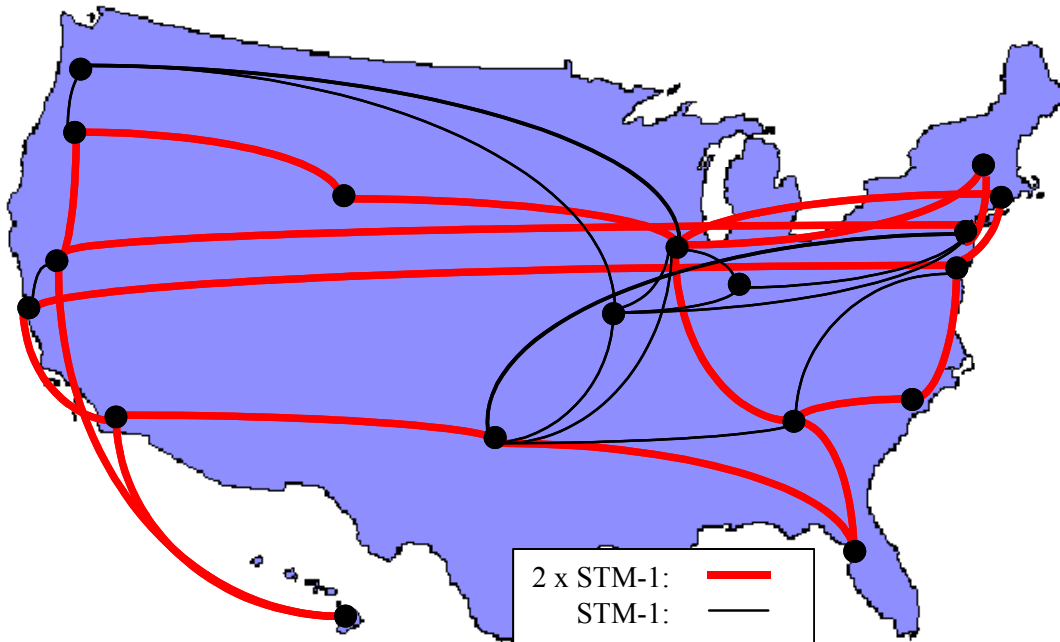
Optimized load balancing solution are computed by the TE tool

- for the worst cases of single link failures, which cause the largest reduction of allowable throughput including those, whose throughput is no more than 10% better than the worst case(s) in combination with
 - all link upgrades
 - Score function for measuring the improvement of a link upgrade with regard to the considered worst failures
- ⇒ The link yielding the highest score is to be upgraded next

Capacity assignment with regard to single link failures

➤ Example: Logical Sprint IP backbone topology [Fraleigh et al. 2003]

➤ Allowable effective mean load: $L = \lambda_{\max}^{def} \sum_{j,k=1}^N d_{\min}(j,k)t_{jk} / \sum_{e \in E} cap(e) \approx 66\%$



$d_{\min}(j,k)$: number of hops on the shortest $j \rightarrow k$ path

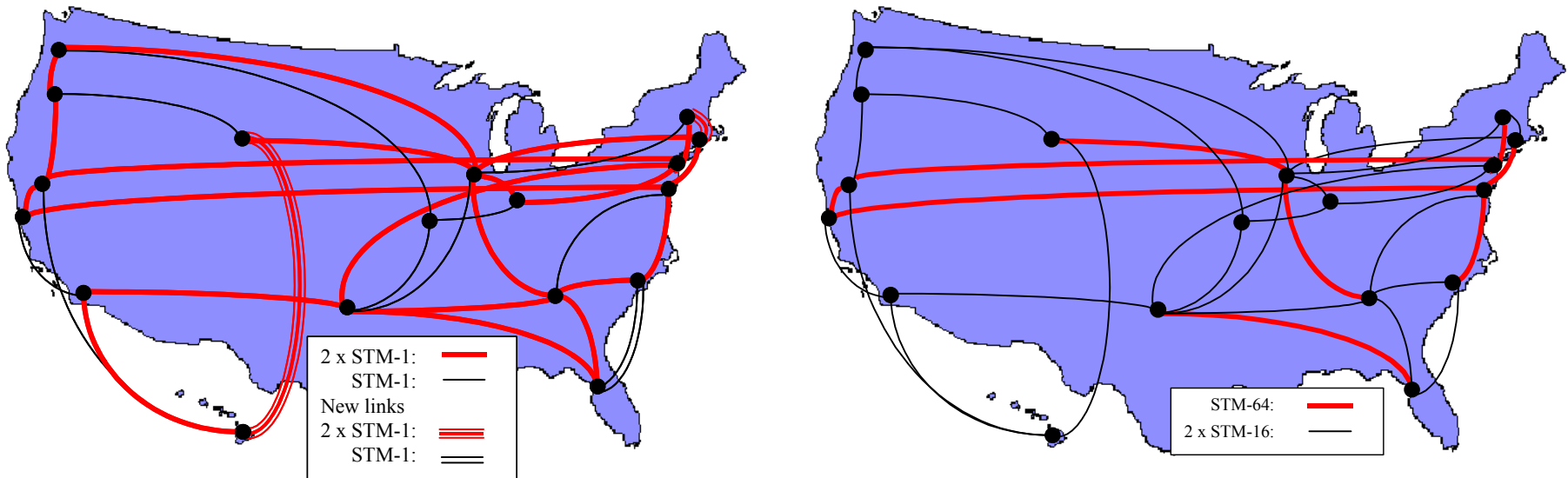
t_{jk} : traffic demand $j \rightarrow k$

$\lambda_{\max} T$: maximum allowable traffic throughput with TE optimization

$cap(e)$: capacity of link e

Modified topology with connectivity degree of 3

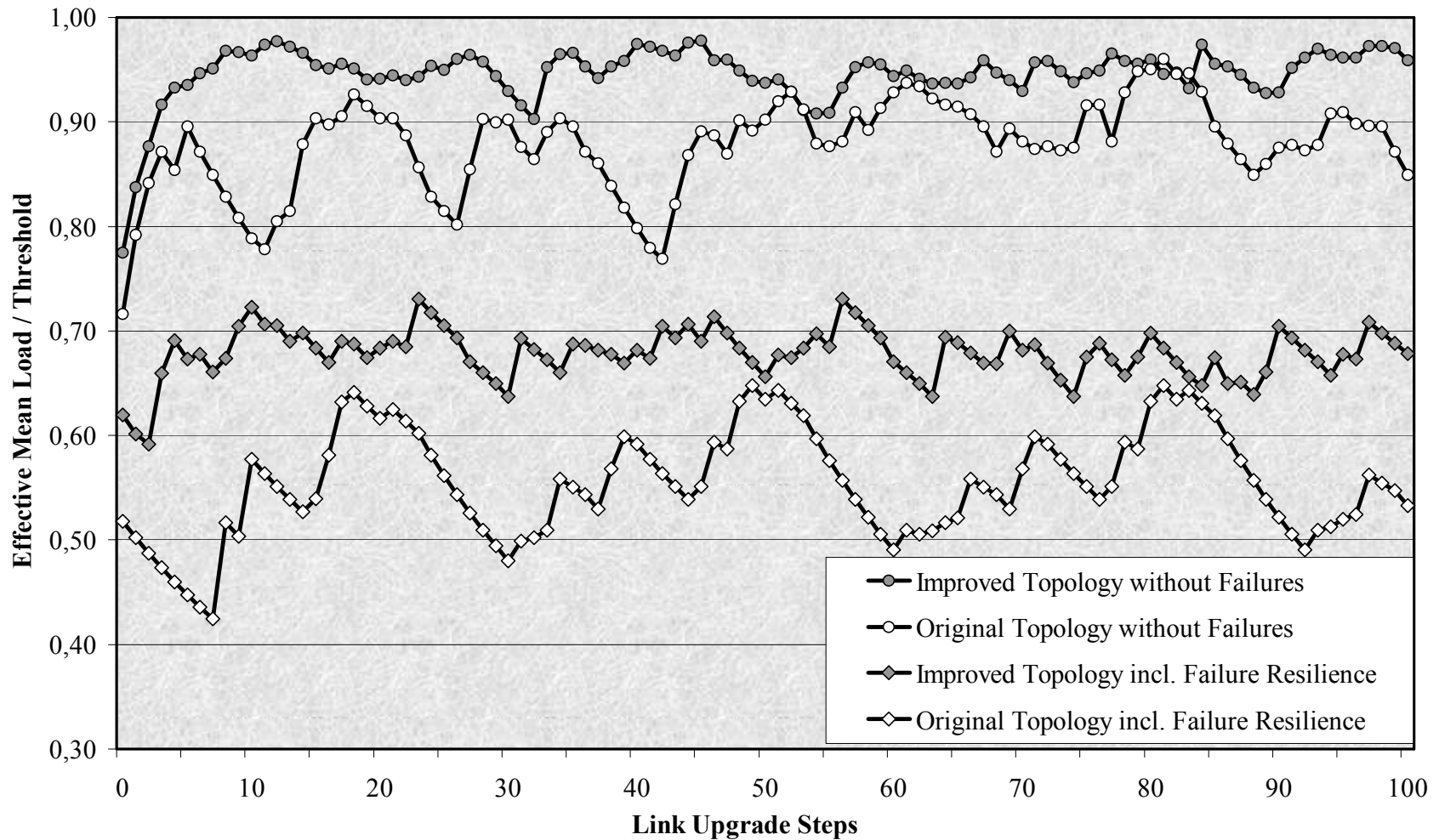
Different capacity assignments achieving 73% effective mean load encountered after 23 and after 174 steps of a link upgrade process



The long term effective mean load is

- 56% in the original topology
- 68% in the modified topology

Ressource utilization during a link by link capacity expansion process for increasing traffic



Long term mean effective load during a link upgrade process

Percentage relative to threshold regarding QoS demands	Original topology	Modified topology (connectivity degree 3)
Traffic Engineering (TE) for normal operation	89 %	94 %
Shortest path routing (SPF) for normal operation	72 %	72 %
Traffic Engineering (TE) including failure resilience	56 %	68 %
Shortest path routing (SPF) including failure resilience	42 %	47 %

Conclusions & Outlook

- Load balancing yields about 20-30% more throughput during upgrade processes for increasing traffic in normal operation
- The traffic engineering gain compared to shortest path routing
 - improves when failure resilience is included
 - improves for optimized topologies (higher connectivity)
- Special upgrade strategies are required for TE; computation of alternative upgrading decisions with the help of TE-tools
- For further study: Upgrading strategies;
Cost functions extending the goal of resource utilization in optimization of topology and capacity assignment