



What's after the Internet?

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Communication today

- **Multiple communication technologies co-exist**
- **Wireless is everywhere, bandwidth will be soon higher than wired**
- **All devices can communicate**
 - multiple network interfaces (comes for free)
 - tons of cheap memory,
 - maybe soon tons of batteries
- **Communication services available through multiple technologies**

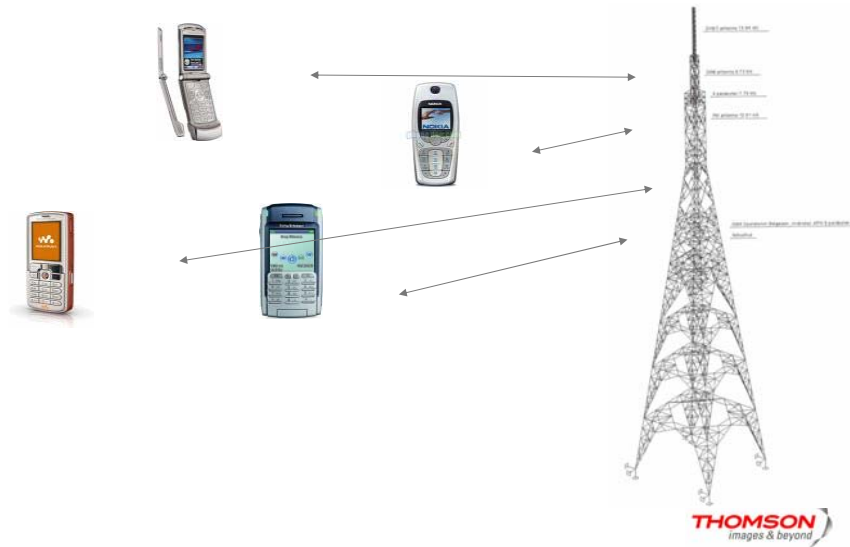


What does not work

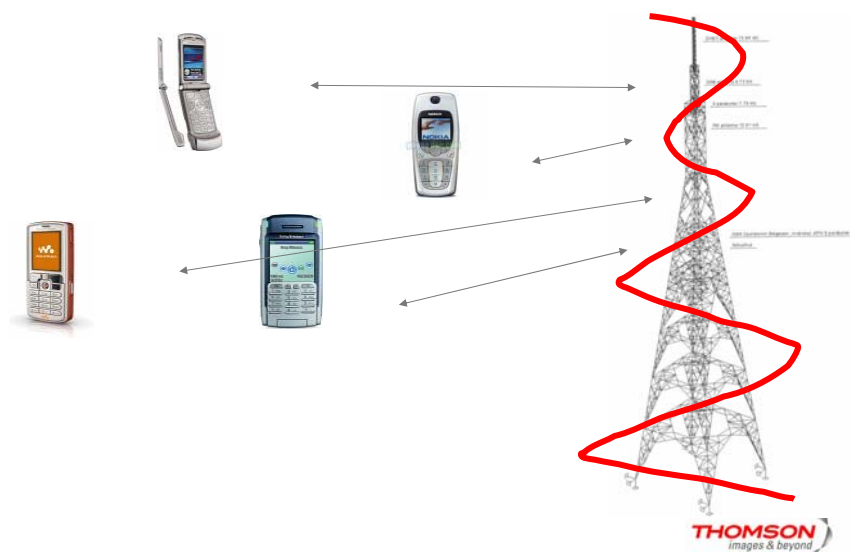
- **The Internet**
 - Not robust
 - Not secure
 - Not capable of integrating wireless (TCP)
 - Not capable of integrating mobility (addressing)
 - Infrastructure based (centralized)
- **The Internet remains the best file transfer network**



Infrastructure based services



Infrastructure based services

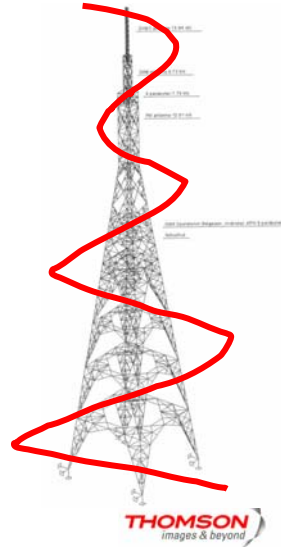


Infrastructure based services

But mobile phones and PDAs are still working....



No infrastructure
Every terminal is also a router
Mobility change topology



No alternative to global services

Today



OR ...



Tomorrow



A world of content

- **Large diversity**
- **Easy to create and distribute**
 - Internet
- **Access to content remains problematic**
 - A dedicated infrastructure per service/content
 - Not interoperable
 - No infrastructure, no service
 - No dedicated terminal, no content





Summary

- **Technology convergence will not happen!**
 - Infrastructures exist and are performant
 - **Content must be available**
 - Everywhere and anytime
 - On all terminals
 - With an « acceptable » quality
- ⇒ **Convergence is happening at the service level**
(triple-play, TV on cell phone, etc.)



A new era for communication

- **New applications, new usages (local & global)**
 - Service personalization (health, tourism)
 - Security (natural disasters)
 - Ease of commercial transactions
 - Legacy applications everywhere
- **New business models**
- **Difficult research problems**
 - Distributed architectures, self-configuration
 - Peer-to-peer search engines
 - Authentication, confidentiality, DRM



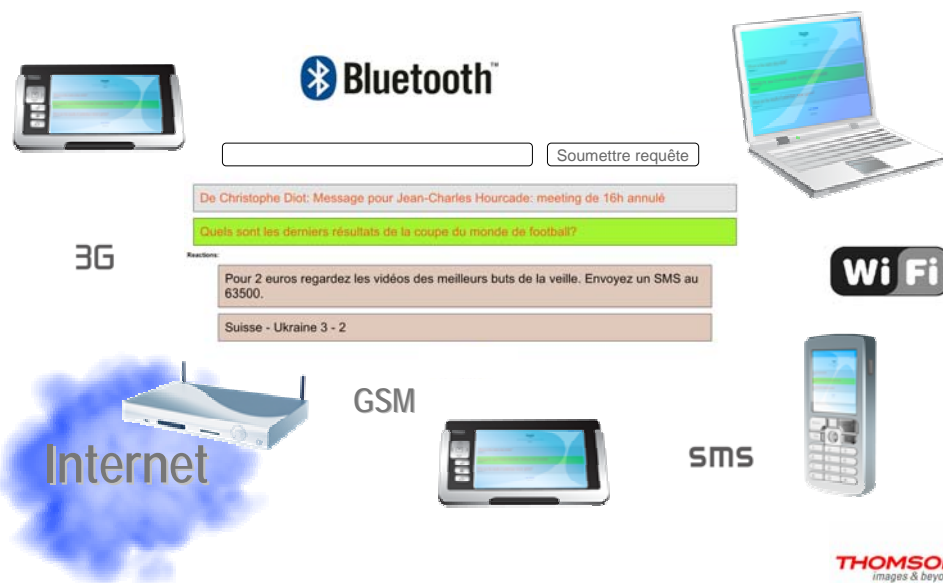
A new communication paradigm

● Pocket Switched Networks

- Any network capable device can store and forward content
- Exchange content using device/user mobility
- Local peer-to-peer communication
- Access to legacy services through gateways



How would it work?



Pocket Switched Networks

- **Provide a useful set of services in the absence of a fixed infrastructure and e2e connectivity**
- **Cross devices, cross network technologies**
- **Exploit features of the problem space**
 - node mobility for message delivery
 - build communities
 - distributed, intuitive authentication
- **Integration with legacy systems**
 - email delivery
 - web requests



Assumptions (1)

- **Users carry devices with connectivity**
 - Bluetooth, 802.11, Ethernet, cradle, etc.
- **Storage is not a problem**
 - storage density doubling every year
- **Battery is more of a problem**
 - heuristics to determine how scarce battery is
 - scale Huggle operations appropriately
 - allow user control



Assumptions (2)

- Local (e.g. Bluetooth, WiFi) and global (e.g. Internet, GPRS) connectivity are *transfer opportunities*
- Use application-layer data instead of (address, port) for forwarding
- No centralized service available (naming, security, etc.)



Challenges

- **Exploit massive aggregate bandwidth**
 - Devices with local connectivity
 - Make use of MBs of local storage
 - Heterogeneous network types
- **Distributed naming**
- **Nodes need to “locate” themselves and their neighbours**
- **Forwarding decision**
- **Security, trust and reputation**
- **Paid to user acceptance, usability, and privacy**



Community networks

- **Recently, many applications have relied on communities sharing a common interest/goal:**
 - Overlay networks, VPNs, etc.
 - File sharing P2P networks
 - Ad-hoc networks
- **Communities may be transient (concert attendees) or long-lived (interest groups)**
- **Users may be in multiple communities at any time, and change communities over time**
- **Issues with naming, trustability, security, incentive to cooperate**



Small world forwarding

- **State information are flooded**
- **No routes, no routing**
- **Network Coding ???**
- **Application forwarding based on neighbourhood status and history:**
 - Find the next hop that has the highest probability to deliver
 - Avoid flooding
 - Data aging



Localization

- **Node localization is important**
 - Neighbourhood discovery/community creation
 - Location-aware applications
 - Trust and security
- **Various options**
 - Relative to other nodes
 - Absolute (e.g. GPS)
 - Based on some external service (c.f. PlaceLab use of base stations)



Trust and Security

- **Human in the loop for bootstrapping trust**
- **New paradigms for trust establishment:**
 - history of past interactions
 - incentives based on virtual cash
 - Recommendation schemes
- **Opportunistic forwarding vs. cryptography: combine encryption with network coding**
- **Overall privacy preserving mechanisms**



Usability and User Involvement

- **Usability will require informing the user about network state issues – but how?**
 - Spinning globe, signal strength bars, greyed-out names
- **When is user involvement required?**
 - Most decisions must be made by devices
 - Some decisions may have to be left to user
- **What incentive to cooperate?**



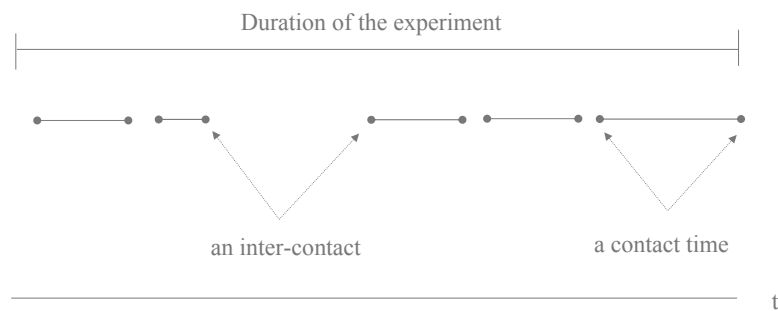
Traces driven design of forwarding

- **Bluetooth (iMotes)**
 - Cambridge local
 - Infocom 05
 - Hong-Kong
- **WiFi (Laptops)**
 - Dartmouth
 - San Diego
- **GSM & Bluetooth (cell phones)**
 - Reality Mining (MIT)

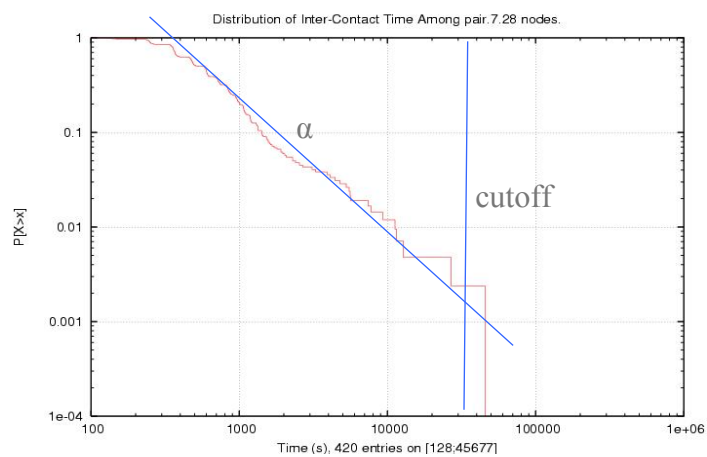


What we measure

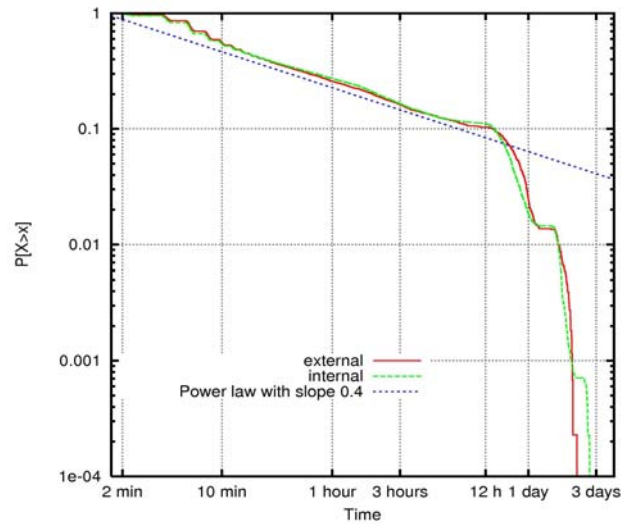
- For a given pairs of devices:
 - contact times and inter-contact times.



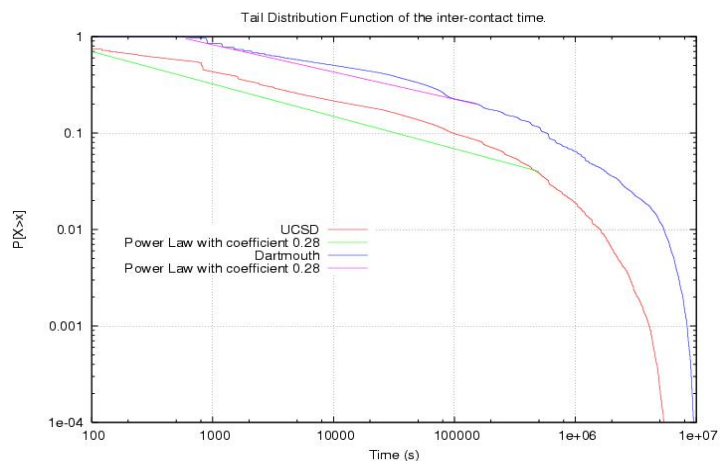
Example: a typical pair



Inter-contact time – iMote experiments



Inter-contact time -- WiFi experiments



Summary of observations

- Inter-contact time follows an approximate power-law shape in all experiments in the [10mns, 12 hours] range.
- $\alpha < 1$ most of the time (very heavily tailed).
- Some variability with the time of day, or among pairs.



Impact of α on convergence

- **For $\alpha > 2$**
Any stateless algorithm achieves a finite expected delay.
- **For $\alpha \geq \frac{m+1}{m}$ and $\# \{ \text{nodes} \} \geq 2m$:**
There exist a forwarding algorithm with m copies and a finite expected delay.
- **For $\alpha < 1$**
No stateless algorithm (even flooding) achieve a bounded delay (Orey's theorem).

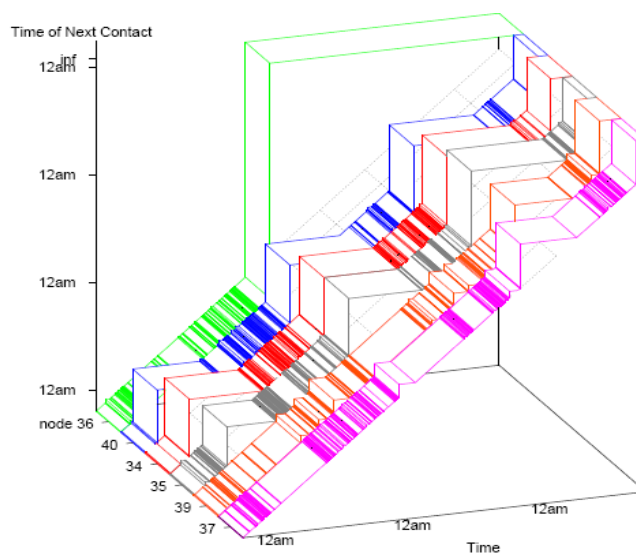


Consequences on mobile networking

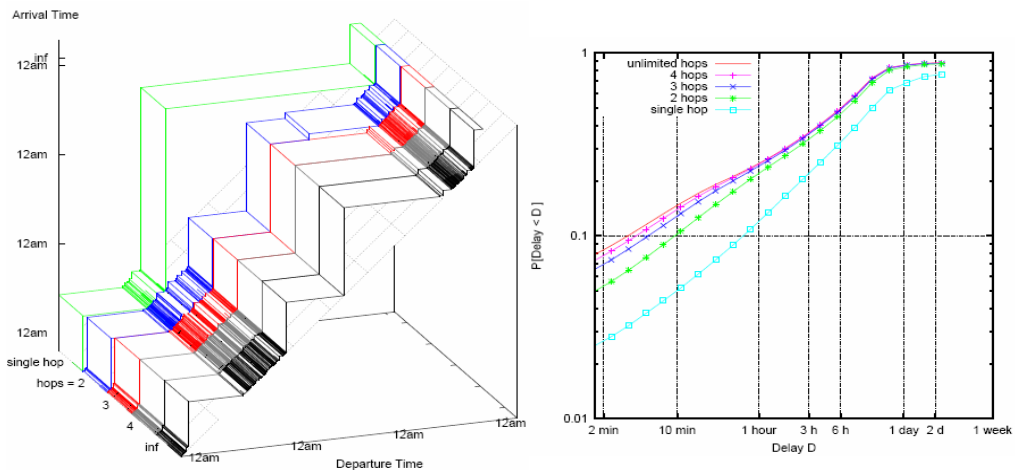
- **Mobility models needs to be redesigned**
 - Exponential decay of inter contact is wrong.
 - Mechanisms tested with that model need to be analyzed with new mobility assumptions.
- **Forwarding will not be easy**
 - Are there simple forwarding rules?
 - Can we benefit from heterogeneity to forward by communities ?



Mobility increases opportunities



Forwarding in PSN



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On-going work

- **Study the limits of naïve forwarding**
 - 4 hops, no more than 8 copies
- **Are we a small world?**
 - Kleinberg's result would apply
- **Characterization of optimal paths**
- **Communities?**
 - Do they exist
 - Can we use them to forward?
 - How?

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URLs

- www.hagggleproject.net
- www.hagggleproject.com
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